8086 Microprocessor

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First Generation
Between 1971 – 1973
PMOS technology, non compatible with TTL
4 bit processors ⇒ 16 pins
8 and 16 bit processors ⇒ 40 pins
Due to limitations of pins, signals are multiplexed

Second Generation
During 1973
NMOS technology ⇒ Faster speed, Higher density, Compatible with TTL
4 / 8/ 16 bit processors ⇒ 40 pins
Ability to address large memory spaces and I/O ports
Greater number of levels of subroutine nesting
Better interrupt handling capabilities

Intel 8085 (8 bit processor)

Third Generation
During 1978
HMOS technology ⇒ Faster speed, Higher packing density
16 bit processors ⇒ 40/ 48/ 64 pins
Easier to program
Dynamically relatable programs
Processor has multiply/ divide arithmetic hardware
More powerful interrupt handling capabilities
Flexible I/O port addressing

Intel 8086 (16 bit processor)

Fourth Generation
During 1980s
Low power version of HMOS technology (HCMOS)
32 bit processors
Physical memory space 2^{24} bytes = 16 Mb
Virtual memory space 2^{40} bytes = 1 Tb
Floating point hardware
Supports increased number of addressing modes

Intel 80386

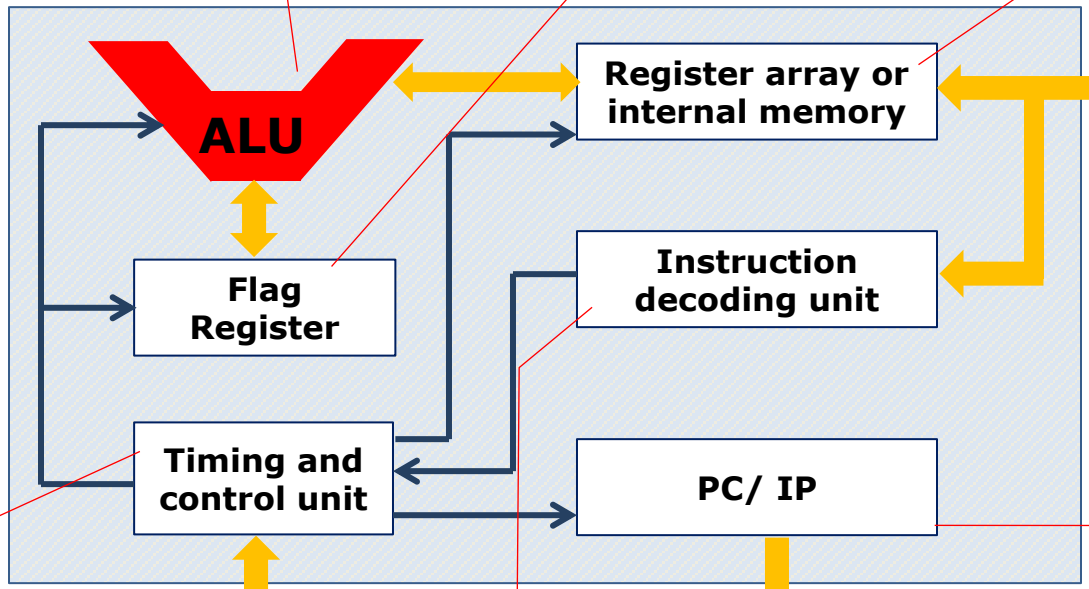
Fifth Generation Pentium

Functional blocks

Computational Unit; performs arithmetic and logic operations

Various conditions of the results are stored as status bits called flags in flag register

Internal storage of data



Generates the address of the instructions to be fetched from the memory and send through address bus to the memory

Generates control signals for internal and external operations of the microprocessor

Decodes instructions; sends information to the timing and control unit

Overview

First 16- bit processor released by INTEL in the year 1978

Originally HMOS, now manufactured using HMOS III technique

Approximately 29, 000 transistors, 40 pin DIP, 5V supply

Does not have internal clock; external asymmetric clock source with 33% duty cycle

20-bit address to access memory ⇒ can address up to 2^{20} = 1 megabytes of memory space.

Addressable memory space is organized in to two banks of 512 kb each; Even (or lower) bank and Odd (or higher) bank. Address line A_0 is used to select even bank and control signal $\overline{BHE}$ is used to access odd bank

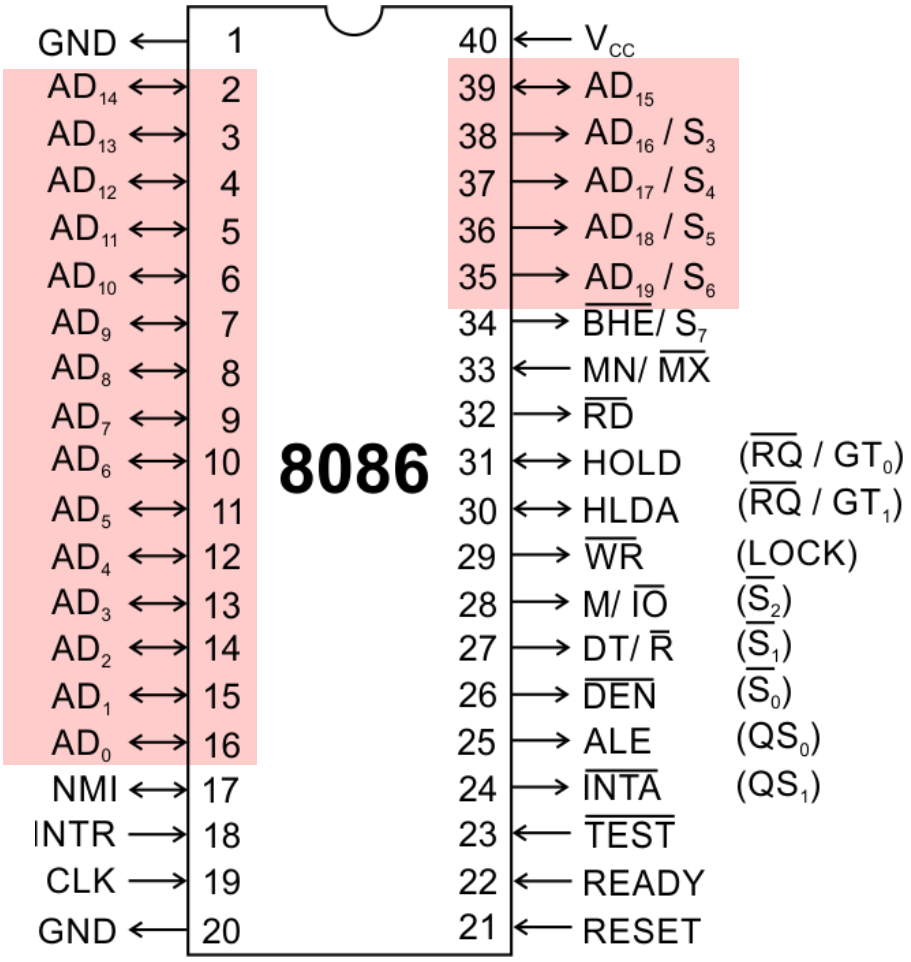
Uses a separate 16 bit address for I/O mapped devices ⇒ can generate 2^{16} = 64 k addresses.

Operates in two modes: minimum mode and maximum mode, decided by the signal at MN and $\overline{MX}$ pins.

Pins and signals

Pins and Signals

Common signals



AD_0-AD_{15} (Bidirectional)

Address/Data bus

Low order address bus; these are multiplexed with data.

When AD lines are used to transmit memory address the symbol A is used instead of AD, for example A_0-A_{15}.

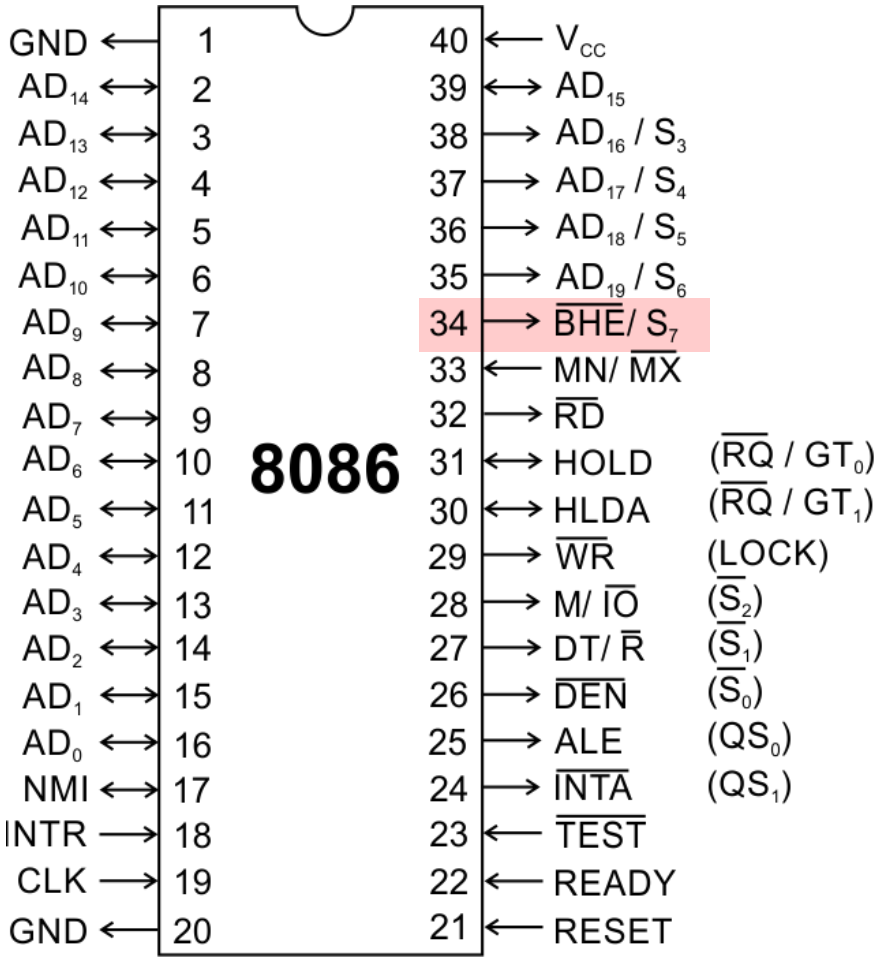
When data are transmitted over AD lines the symbol D is used in place of AD, for example D_0-D_7, D_8-D_{15} or D_0-D_{15}.

A_{16}/S_3, A_{17}/S_4, A_{18}/S_5, A_{19}/S_6

High order address bus. These are multiplexed with status signals

Pins and Signals

Common signals



BHE (Active Low)/S_7 (Output)

Bus High Enable/Status

It is used to enable data onto the most significant half of data bus, D_8-D_{15}. 8-bit device connected to upper half of the data bus use BHE (Active Low) signal. It is multiplexed with status signal S_7.

MN/ MX

MINIMUM / MAXIMUM

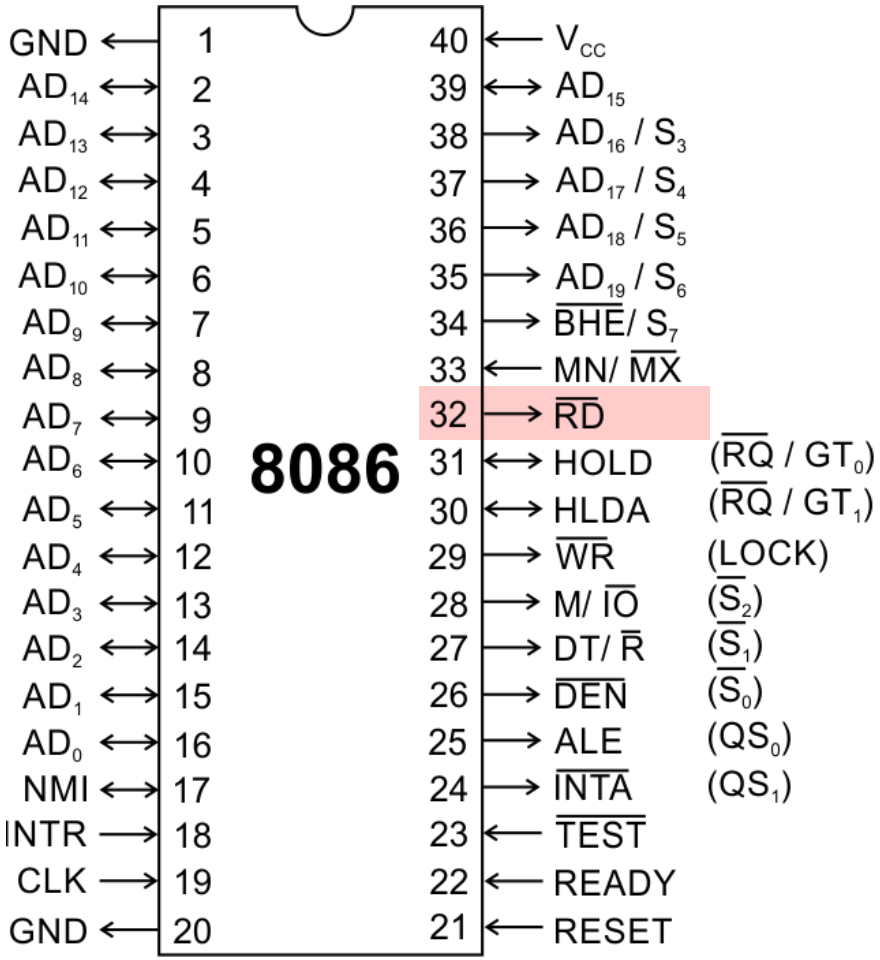
This pin signal indicates what mode the processor is to operate in.

RD (Read) (Active Low)

The signal is used for read operation.
It is an output signal.
It is active when low.

Pins and Signals

Common signals



TEST

TEST input is tested by the 'WAIT' instruction.

8086 will enter a wait state after execution of the WAIT instruction and will resume execution only when the TEST is made low by an active hardware.

This is used to synchronize an external activity to the processor internal operation.

READY

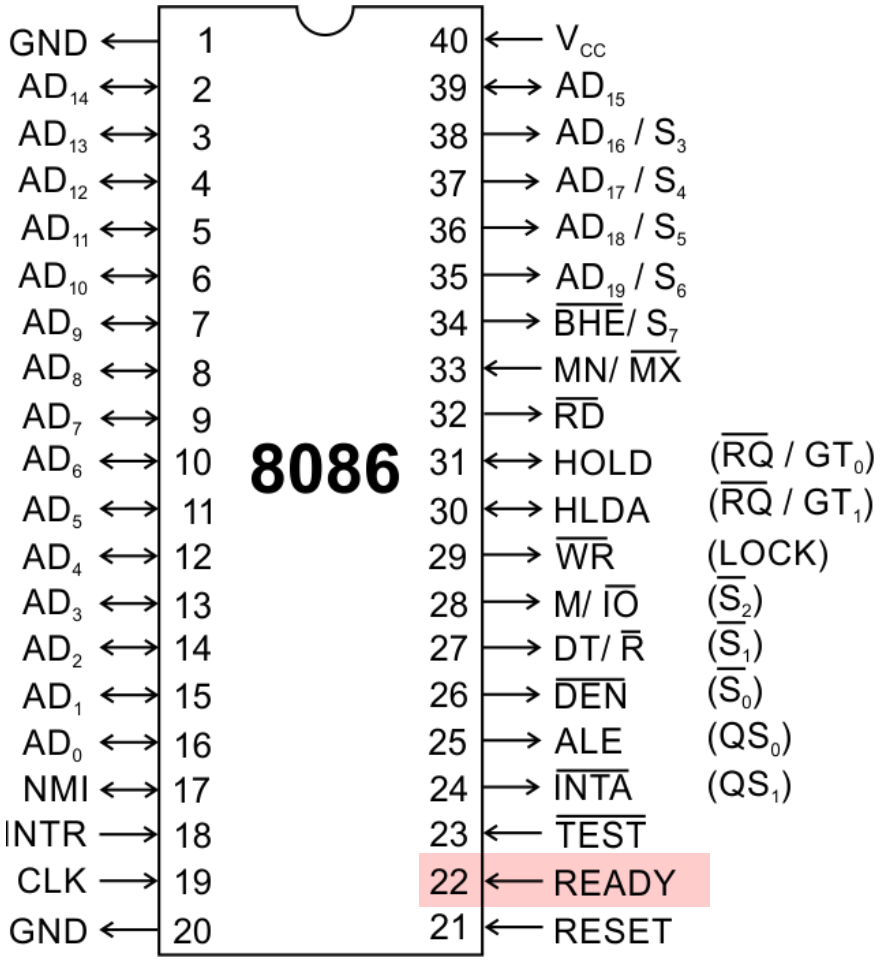
This is the acknowledgement from the slow device or memory that they have completed the data transfer.

The signal made available by the devices is synchronized by the 8284A clock generator to provide ready input to the 8086.

The signal is active high.

Pins and Signals

Common signals



RESET (Input)

Causes the processor to immediately terminate its present activity.

The signal must be active HIGH for at least four clock cycles.

CLK

The clock input provides the basic timing for processor operation and bus control activity. Its an asymmetric square wave with 33% duty cycle.

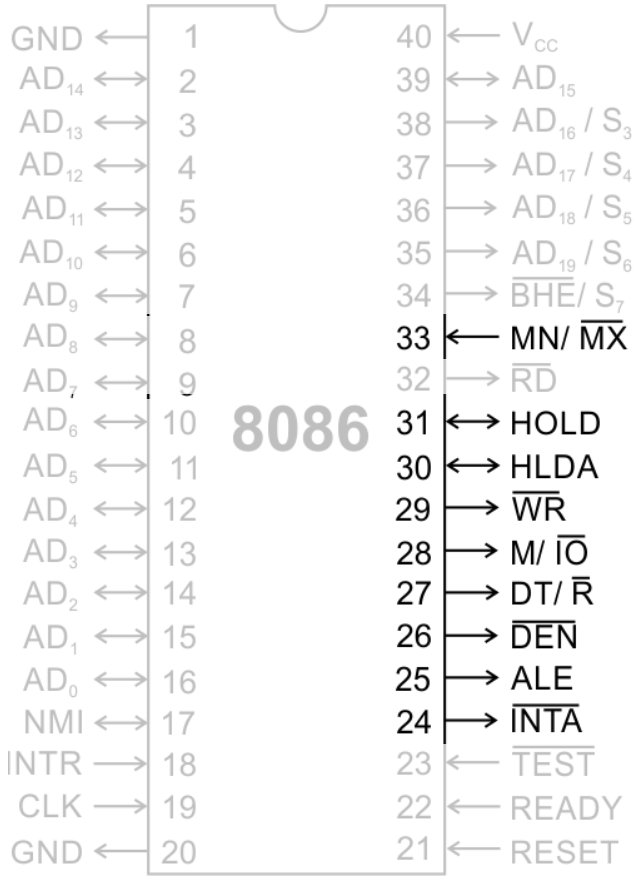
INTR Interrupt Request

This is a triggered input. This is sampled during the last clock cycles of each instruction to determine the availability of the request. If any interrupt request is pending, the processor enters the interrupt acknowledge cycle.

This signal is active high and internally synchronized.

Pins and Signals

Min/ Max Pins



The 8086 microprocessor can work in two modes of operations : Minimum mode and Maximum mode.

In the minimum mode of operation the microprocessor do not associate with any co-processors and can not be used for multiprocessor systems.

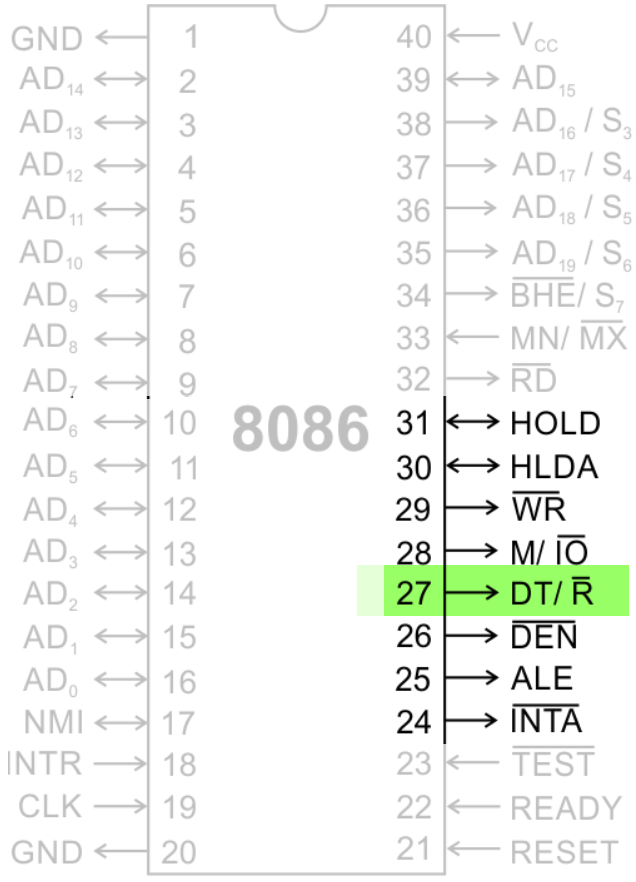
In the maximum mode the 8086 can work in multi-processor or co-processor configuration.

Minimum or maximum mode operations are decided by the pin MN/ MX(Active low).

When this pin is high 8086 operates in minimum mode otherwise it operates in Maximum mode.

Pins and Signals

Minimum mode signals



Pins 24 -31

For minimum mode operation, the MN/ $\overline{MX}$ is tied to VCC (logic high)

8086 itself generates all the bus control signals

DT/$\overline{R}$ — (Data Transmit/ Receive) Output signal from the processor to control the direction of data flow through the data transceivers

$\overline{DEN}$ — (Data Enable) Output signal from the processor used as out put enable for the transceivers

ALE — (Address Latch Enable) Used to demultiplex the address and data lines using external latches

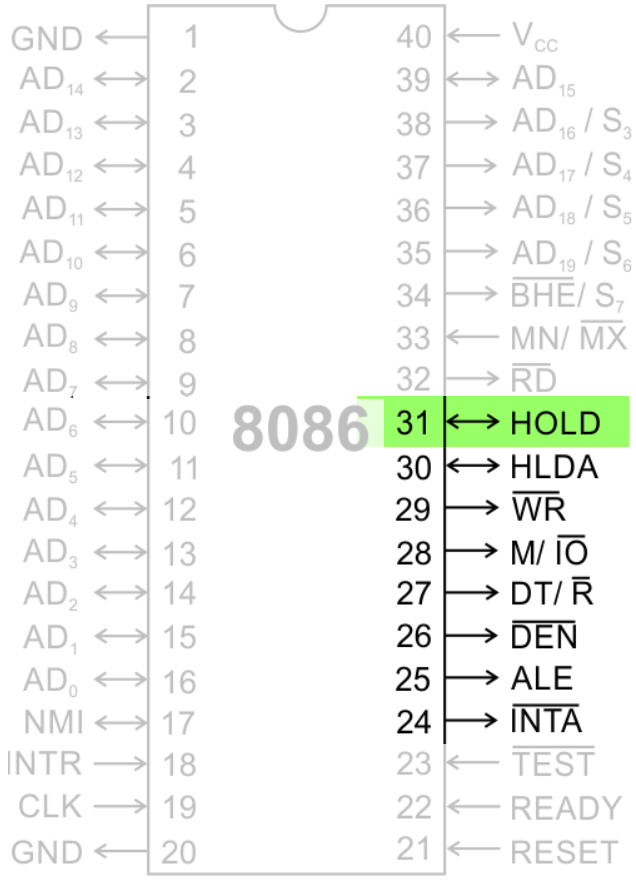
M/$\overline{IO}$ — Used to differentiate memory access and I/O access. For memory reference instructions, it is high. For IN and OUT instructions, it is low.

$\overline{WR}$ — Write control signal; asserted low Whenever processor writes data to memory or I/O port

$\overline{INTA}$ — (Interrupt Acknowledge) When the interrupt request is accepted by the processor, the output is low on this line.

Pins and Signals

Minimum mode signals



Pins 24 -31

For minimum mode operation, the MN/ $\overline{MX}$ is tied to VCC (logic high)

8086 itself generates all the bus control signals

HOLD — **Input signal to the processor form the bus masters as a request to grant the control of the bus.**

Usually used by the DMA controller to get the control of the bus.

HLDA — **(Hold Acknowledge) Acknowledge signal by the processor to the bus master requesting the control of the bus through HOLD.**

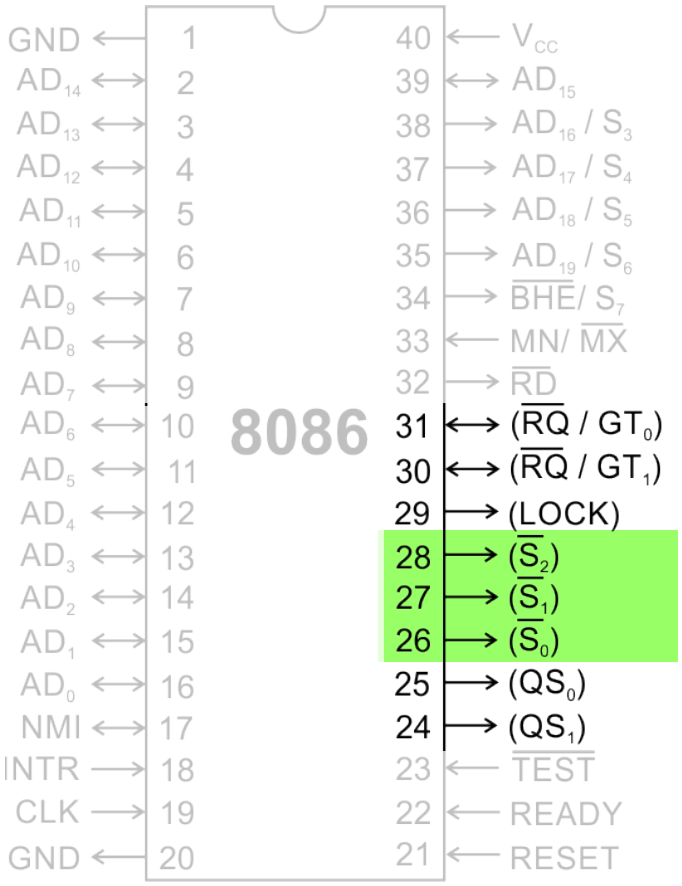
The acknowledge is asserted high, when the processor accepts HOLD.

Pins and Signals

Maximum mode signals

During maximum mode operation, the MN/ $\overline{MX}$ is grounded (logic low)

Pins 24 -31 are reassigned



$\overline{S_0}$, $\overline{S_1}$, $\overline{S_2}$ **Status signals; used by the 8086 bus controller to generate bus timing and control signals. These are decoded as shown.**

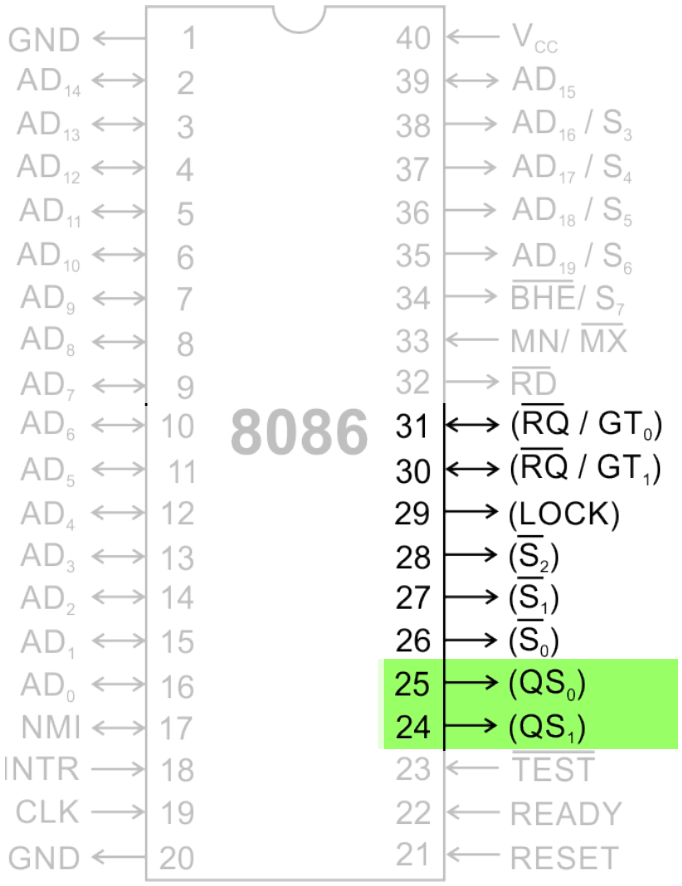
Status Signal			Machine Cycle
$\overline{S}_2$	$\overline{S}_1$	$\overline{S}_0$	
0	0	0	Interrupt acknowledge
0	0	1	Read I/O port
0	1	0	Write I/O port
0	1	1	Halt
1	0	0	Code access
1	0	1	Read memory
1	1	0	Write memory
1	1	1	Passive/Inactive

Pins and Signals

Maximum mode signals

During maximum mode operation, the MN/ $\overline{MX}$ is grounded (logic low)

Pins 24 -31 are reassigned



$\overline{QS_0}$, $\overline{QS_1}$ **(Queue Status) The processor provides the status of queue in these lines.**

The queue status can be used by external device to track the internal status of the queue in 8086.

The output on QS_0 and QS_1 can be interpreted as shown in the table.

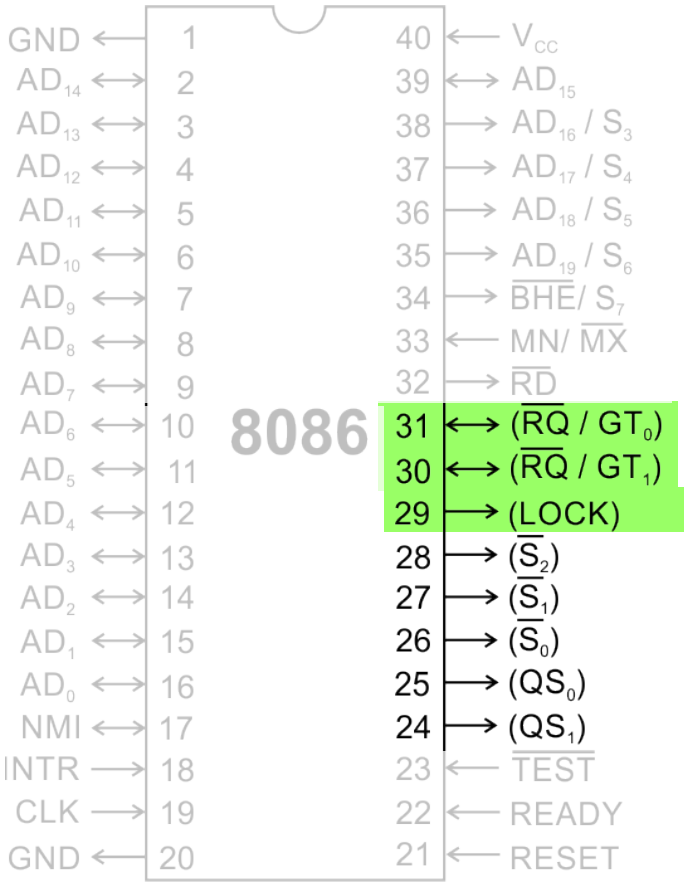
Queue status		Queue operation
QS_1	QS_0	
0	0	No operation
0	1	First byte of an opcode from queue
1	0	Empty the queue
1	1	Subsequent byte from queue

Pins and Signals

Maximum mode signals

During maximum mode operation, the MN/ $\overline{MX}$ is grounded (logic low)

Pins 24 -31 are reassigned



$\overline{RQ}/\ \overline{GT_0}$, $\overline{RQ}/\ \overline{GT_1}$

(**Bus Request/ Bus Grant**) These requests are used by other local bus masters to force the processor to release the local bus at the end of the processor's current bus cycle.

These pins are bidirectional.

The request on $\overline{GT_0}$ will have higher priority than $\overline{GT_1}$

$\overline{LOCK}$

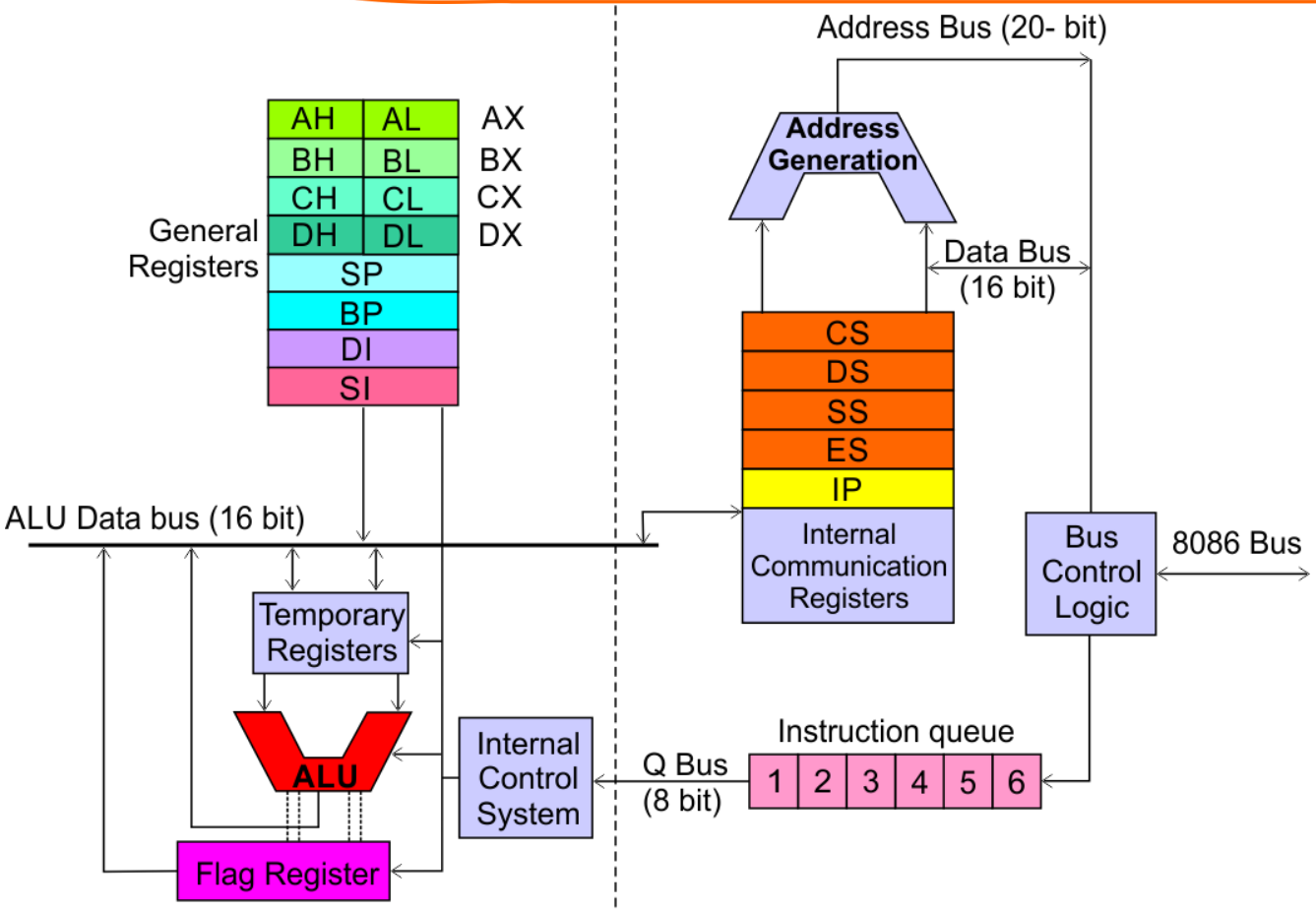
An output signal activated by the LOCK prefix instruction.

Remains active until the completion of the instruction prefixed by LOCK.

The 8086 output low on the $\overline{LOCK}$ pin while executing an instruction prefixed by LOCK to prevent other bus masters from gaining control of the system bus.

Architecture

Architecture



Execution Unit (EU)

EU executes instructions that have already been fetched by the BIU.

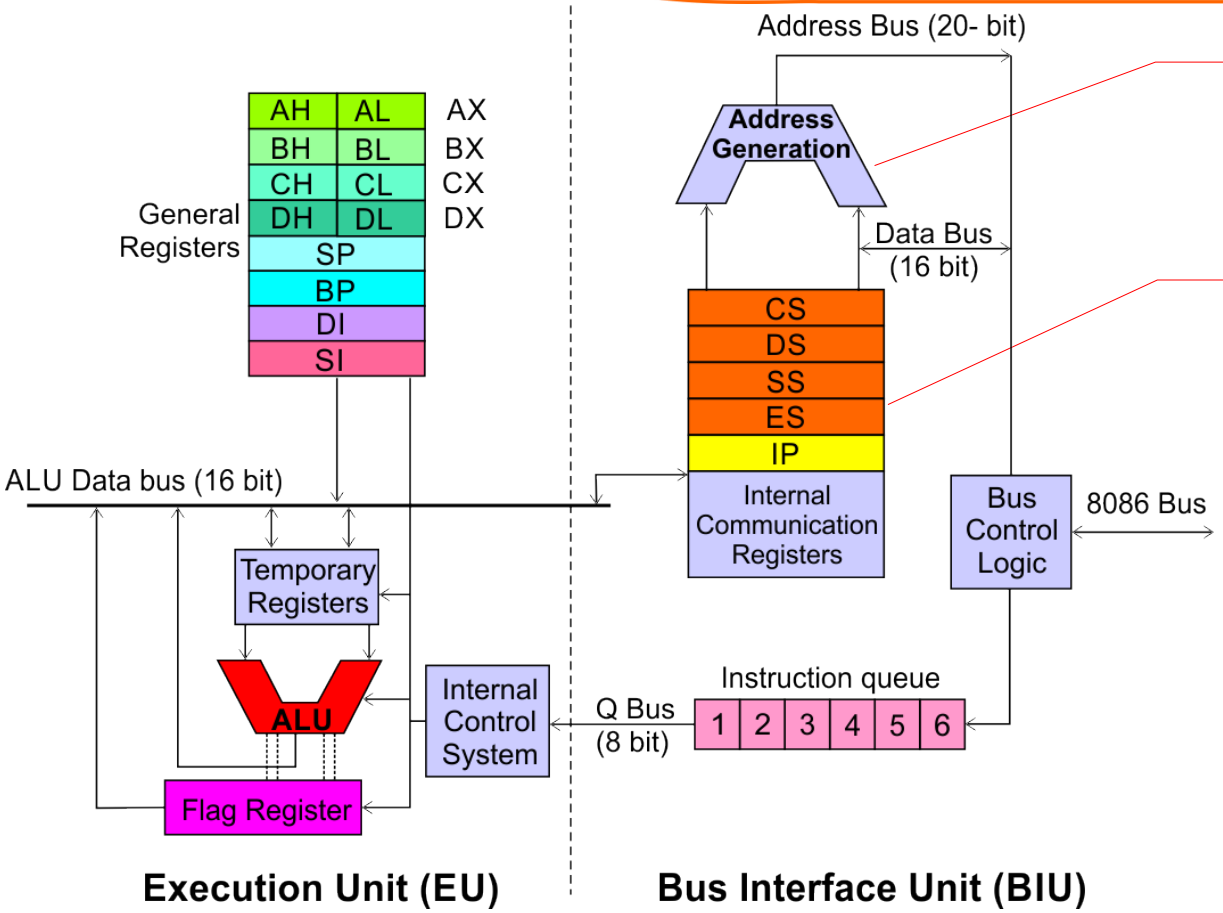
BIU and EU functions separately.

Bus Interface Unit (BIU)

BIU fetches instructions, reads data from memory and I/O ports, writes data to memory and I/ O ports.

Architecture

Bus Interface Unit (BIU)



Dedicated Adder to generate 20 bit address

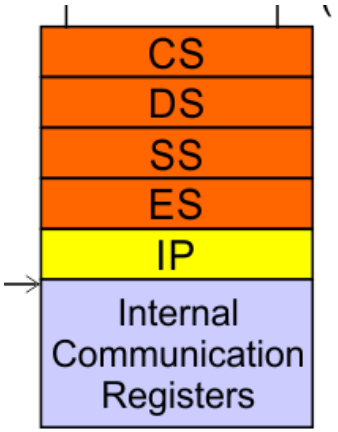
Four 16-bit segment registers

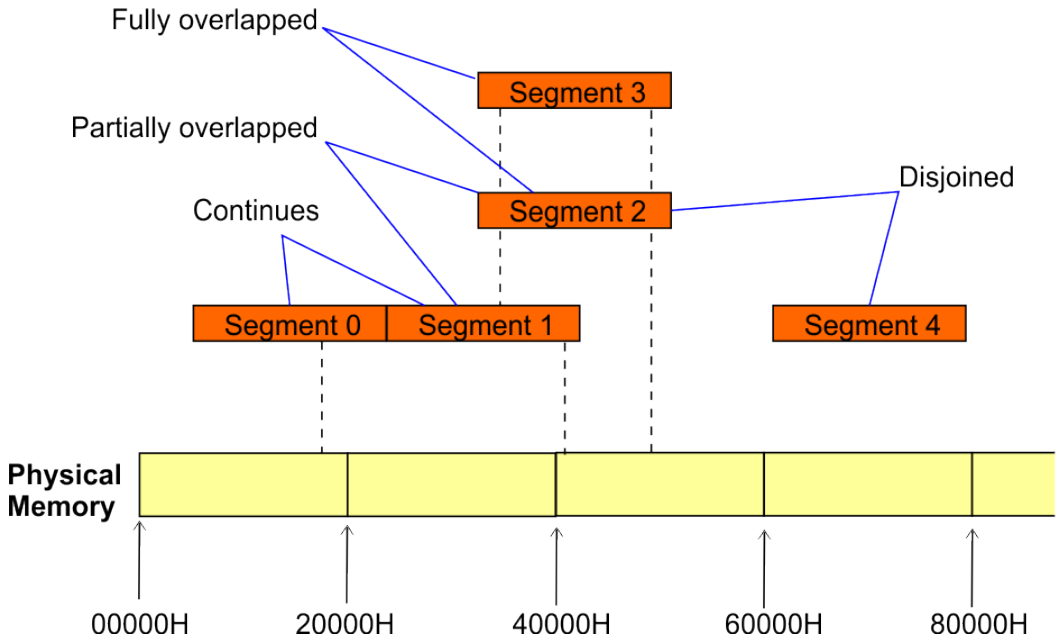
Code Segment (CS)
Data Segment (DS)
Stack Segment (SS)
Extra Segment (ES)

Architecture

Bus Interface Unit (BIU)

Segment Registers



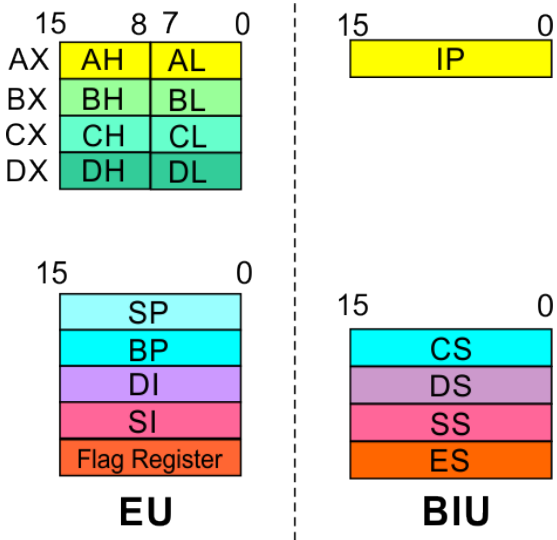


- **8086's 1-megabyte memory is divided into segments of up to 64K bytes each.**
- **The 8086 can directly address four segments (256 K bytes within the 1 M byte of memory) at a particular time.**
- **Programs obtain access to code and data in the segments by changing the segment register content to point to the desired segments.**

Segment Registers

Code Segment Register

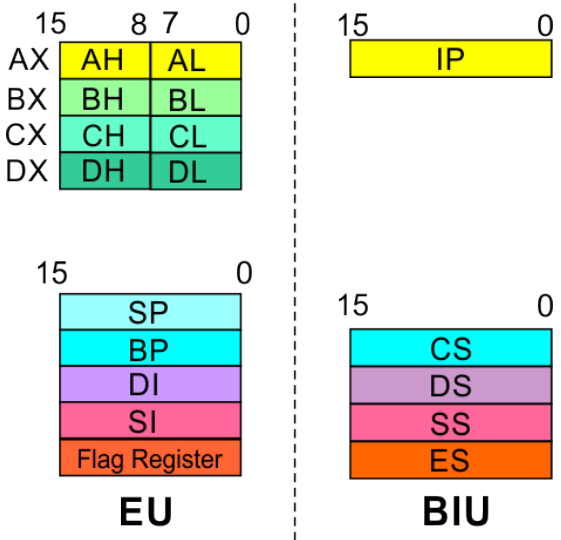
- 16-bit
- CS contains the base or start of the current code segment; IP contains the distance or offset from this address to the next instruction byte to be fetched.
- BIU computes the 20-bit physical address by logically shifting the contents of CS 4-bits to the left and then adding the 16-bit contents of IP.
- That is, all instructions of a program are relative to the contents of the CS register multiplied by 16 and then offset is added provided by the IP.



Segment Registers

Data Segment Register

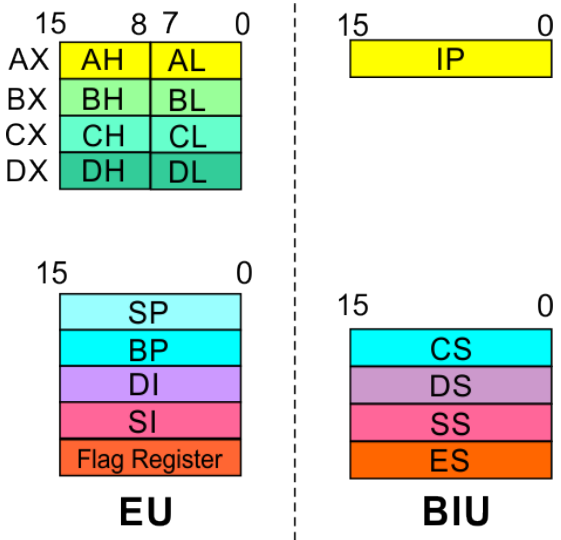
- **16-bit**
- **Points to the current data segment; operands for most instructions are fetched from this segment.**
- **The 16-bit contents of the Source Index (SI) or Destination Index (DI) or a 16-bit displacement are used as offset for computing the 20-bit physical address.**



Segment Registers

Stack Segment Register

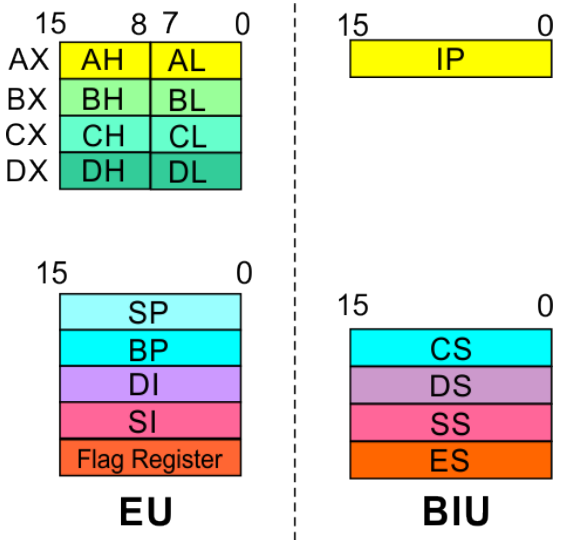
- 16-bit
- Points to the current stack.
- The 20-bit physical stack address is calculated from the Stack Segment (SS) and the Stack Pointer (SP) for stack instructions such as PUSH and POP.
- In based addressing mode, the 20-bit physical stack address is calculated from the Stack segment (SS) and the Base Pointer (BP).



Segment Registers

Extra Segment Register

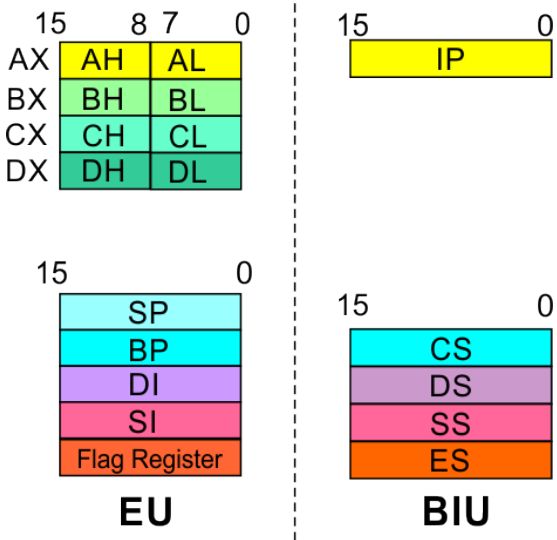
- **16-bit**
- **Points to the extra segment in which data (in excess of 64K pointed to by the DS) is stored.**
- **String instructions use the ES and DI to determine the 20-bit physical address for the destination.**



Segment Registers

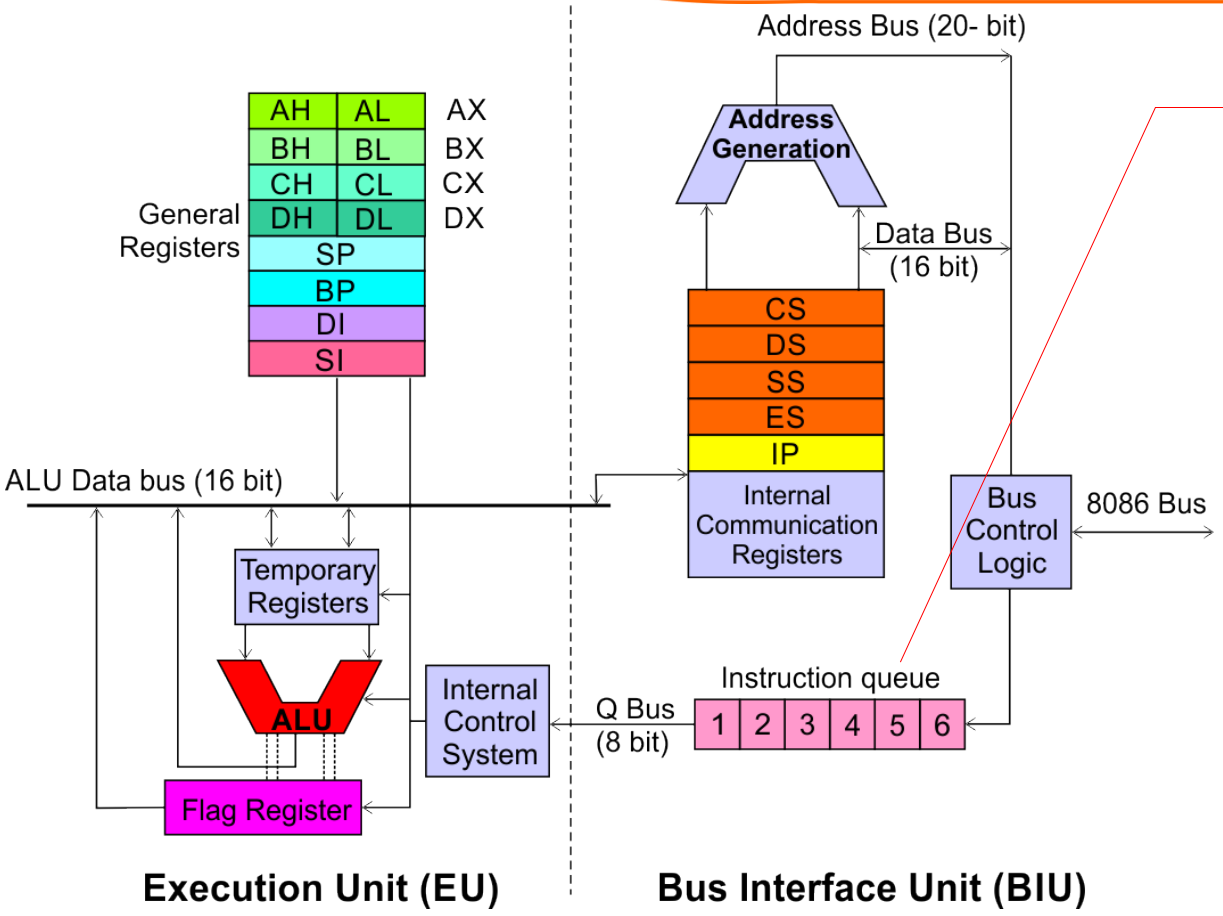
Instruction Pointer

- 16-bit
- Always points to the next instruction to be executed within the currently executing code segment.
- So, this register contains the 16-bit offset address pointing to the next instruction code within the 64Kb of the code segment area.
- Its content is automatically incremented as the execution of the next instruction takes place.



Architecture

Bus Interface Unit (BIU)



Instruction queue

- **A group of First-In-First-Out (FIFO) in which up to 6 bytes of instruction code are pre fetched from the memory ahead of time.**
- **This is done in order to speed up the execution by overlapping instruction fetch with execution.**
- **This mechanism is known as pipelining.**

Architecture

Execution Unit (EU)

EU decodes and executes instructions.

A decoder in the EU control system translates instructions.

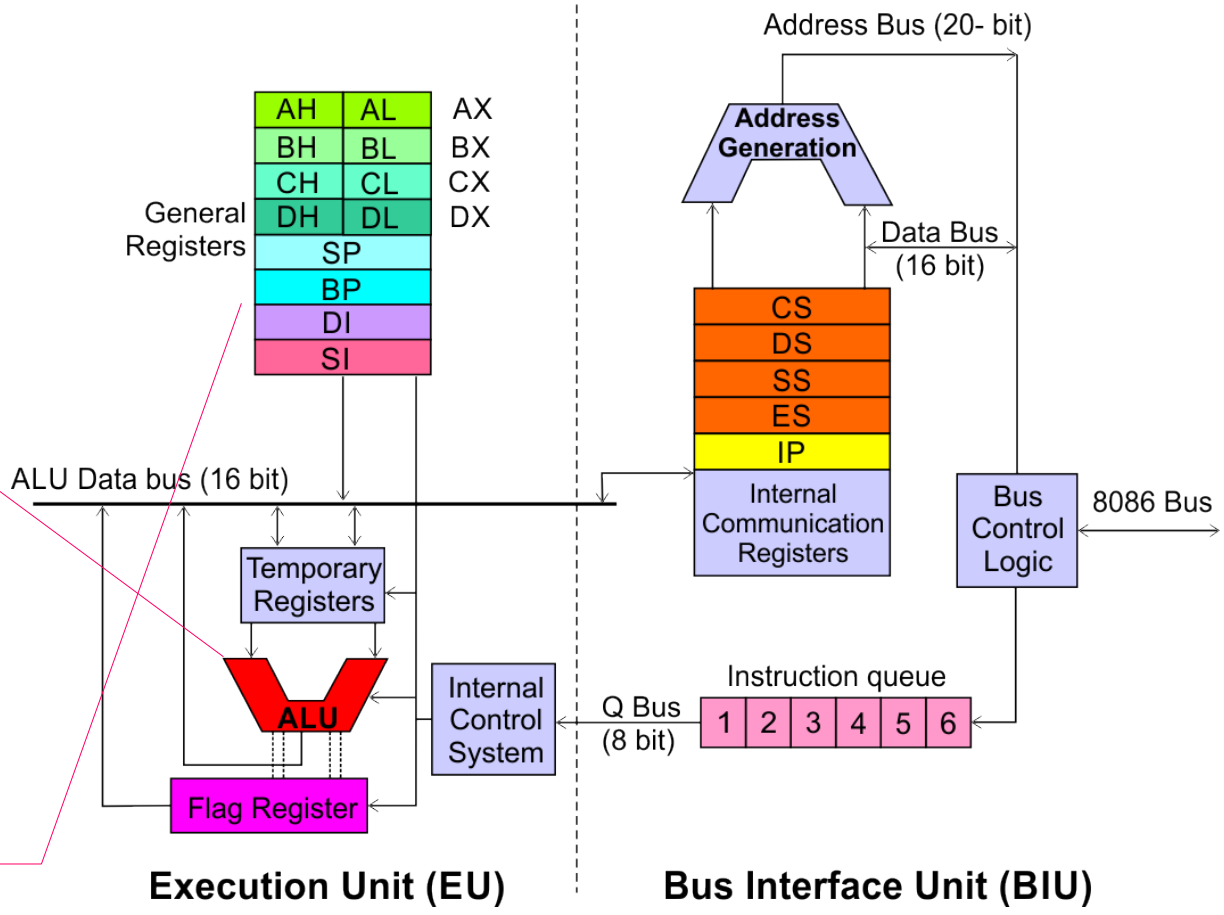
16-bit ALU for performing arithmetic and logic operation

Four general purpose registers(AX, BX, CX, DX);

Pointer registers (Stack Pointer, Base Pointer);

and

Index registers (Source Index, Destination Index) each of 16-bits



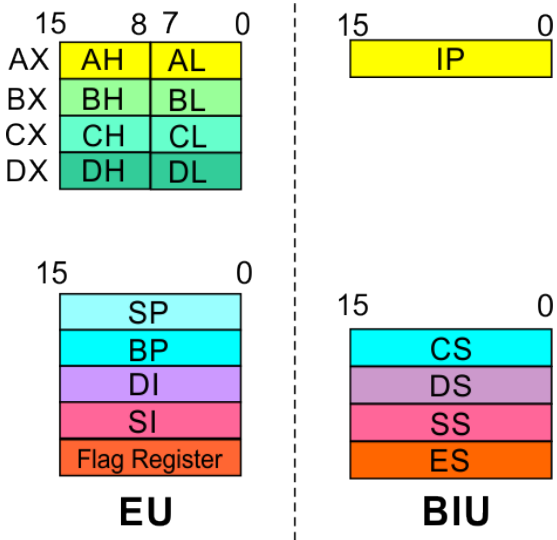
Some of the 16 bit registers can be used as two 8 bit registers as :

AX can be used as AH and AL
BX can be used as BH and BL
CX can be used as CH and CL
DX can be used as DH and DL

EU Registers

Accumulator Register (AX)

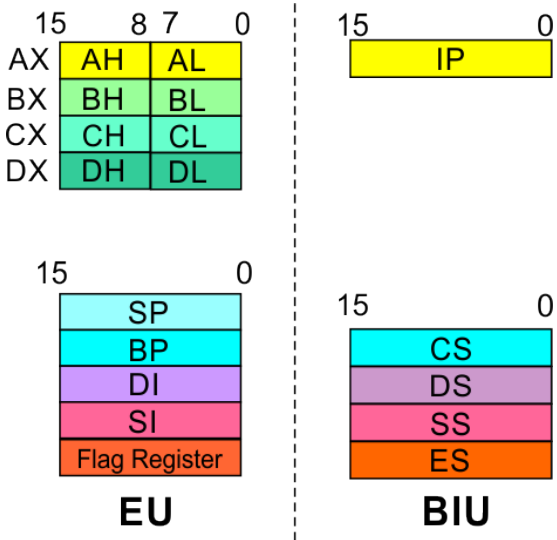
- **Consists of two 8-bit registers AL and AH, which can be combined together and used as a 16-bit register AX.**
- **AL in this case contains the low order byte of the word, and AH contains the high-order byte.**
- **The I/O instructions use the AX or AL for inputting / outputting 16 or 8 bit data to or from an I/O port.**
- **Multiplication and Division instructions also use the AX or AL.**



EU Registers

Base Register (BX)

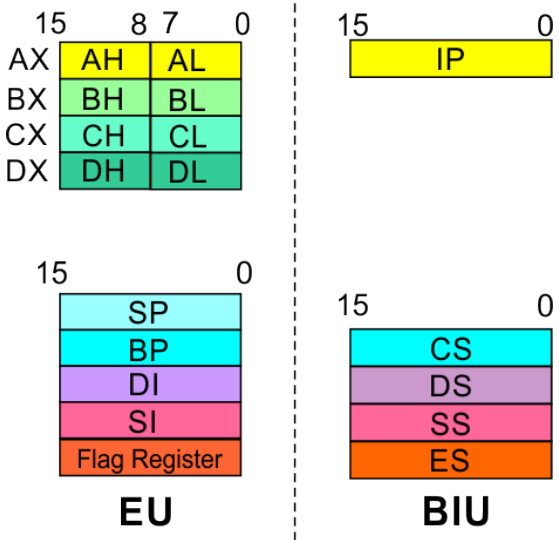
- Consists of two 8-bit registers BL and BH, which can be combined together and used as a 16-bit register BX.
- BL in this case contains the low-order byte of the word, and BH contains the high-order byte.
- This is the only general purpose register whose contents can be used for addressing the 8086 memory.
- All memory references utilizing this register content for addressing use DS as the default segment register.



EU Registers

Counter Register (CX)

- **Consists of two 8-bit registers CL and CH, which can be combined together and used as a 16-bit register CX.**
- **When combined, CL register contains the low order byte of the word, and CH contains the high-order byte.**
- **Instructions such as SHIFT, ROTATE and LOOP use the contents of CX as a counter.**



Example:

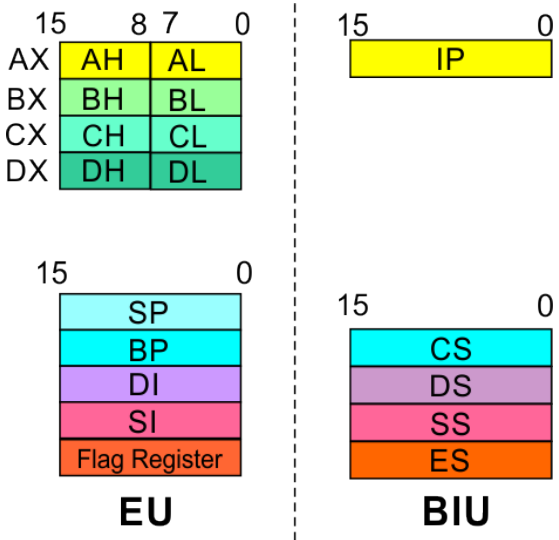
The instruction LOOP START automatically decrements CX by 1 without affecting flags and will check if [CX] = 0.

If it is zero, 8086 executes the next instruction; otherwise the 8086 branches to the label START.

EU Registers

Data Register (DX)

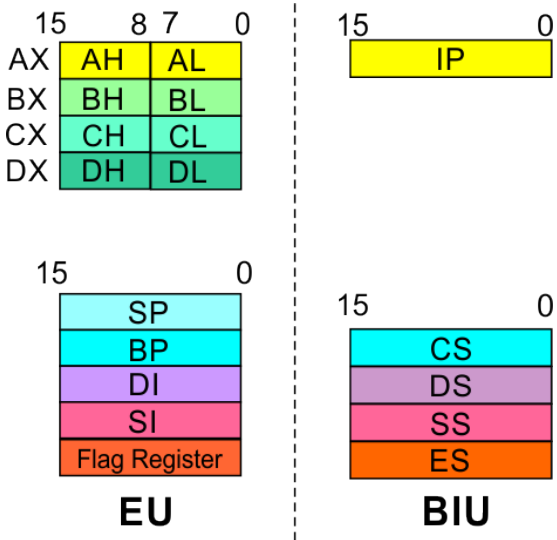
- Consists of two 8-bit registers DL and DH, which can be combined together and used as a 16-bit register DX.
- When combined, DL register contains the low order byte of the word, and DH contains the high-order byte.
- Used to hold the high 16-bit result (data) in 16 X 16 multiplication or the high 16-bit dividend (data) before a 32 ÷ 16 division and the 16-bit reminder after division.



EU Registers

Stack Pointer (SP) and Base Pointer (BP)

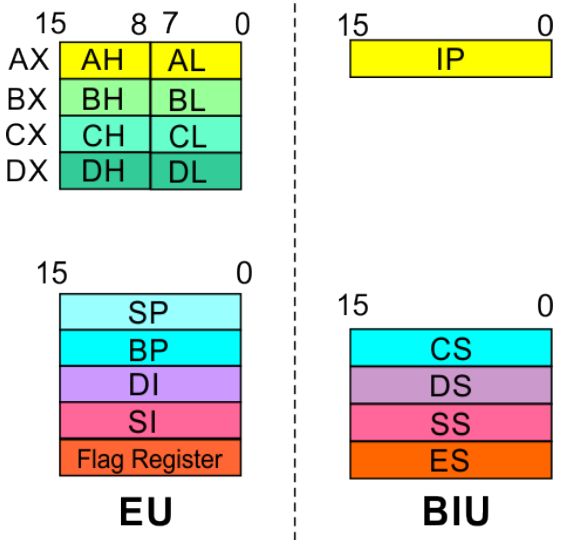
- SP and BP are used to access data in the stack segment.
- SP is used as an offset from the current SS during execution of instructions that involve the stack segment in the external memory.
- SP contents are automatically updated (incremented/ decremented) due to execution of a POP or PUSH instruction.
- BP contains an offset address in the current SS, which is used by instructions utilizing the based addressing mode.



EU Registers

Source Index (SI) and Destination Index (DI)

- **Used in indexed addressing.**
- **Instructions that process data strings use the SI and DI registers together with DS and ES respectively in order to distinguish between the source and destination addresses.**



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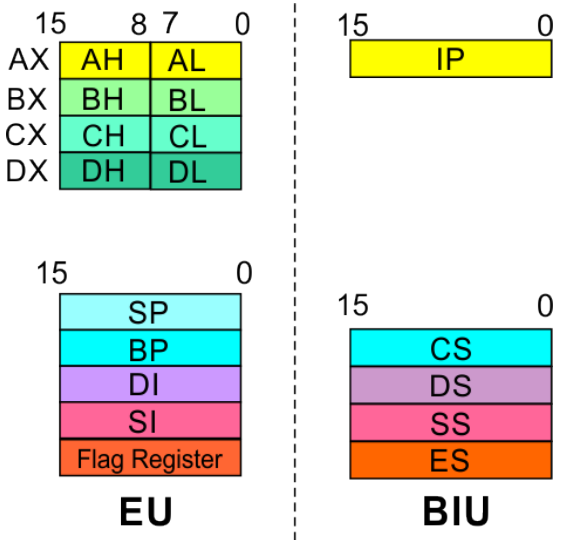


Architecture

Execution Unit (EU)

Flag Register

Auxiliary Carry Flag

This is set, if there is a carry from the lowest nibble, i.e, bit three during addition, or borrow for the lowest nibble, i.e, bit three, during subtraction.

Carry Flag

This flag is set, when there is a carry out of MSB in case of addition or a borrow in case of subtraction.

Sign Flag

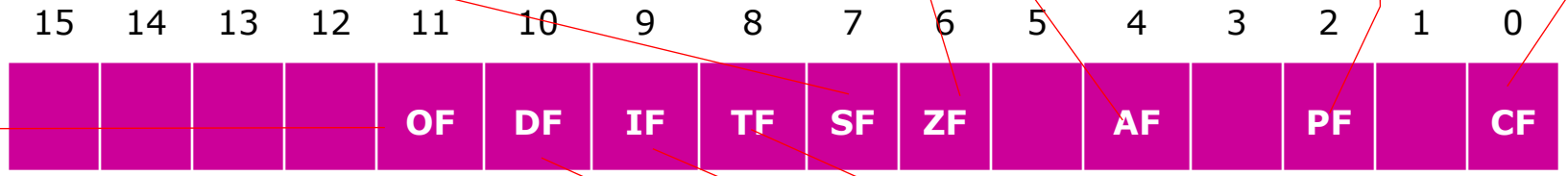
This flag is set, when the result of any computation is negative

Zero Flag

This flag is set, if the result of the computation or comparison performed by an instruction is zero

Parity Flag

This flag is set to 1, if the lower byte of the result contains even number of 1's ; for odd number of 1's set to zero.



Over flow Flag

This flag is set, if an overflow occurs, i.e, if the result of a signed operation is large enough to accommodate in a destination register. The result is of more than 7-bits in size in case of 8-bit signed operation and more than 15-bits in size in case of 16-bit sign operations, then the overflow will be set.

Tarp Flag

If this flag is set, the processor enters the single step execution mode by generating internal interrupts after the execution of each instruction

Direction Flag

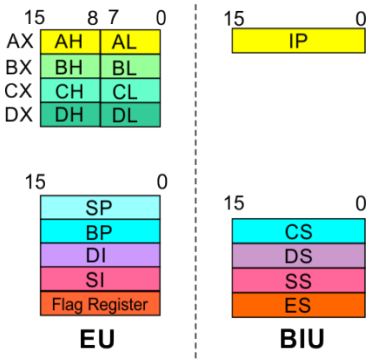
This is used by string manipulation instructions. If this flag bit is '0', the string is processed beginning from the lowest address to the highest address, i.e., auto incrementing mode. Otherwise, the string is processed from the highest address towards the lowest address, i.e., auto incrementing mode.

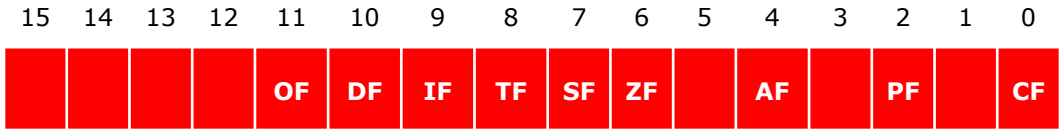
Interrupt Flag

Causes the 8086 to recognize external mask interrupts; clearing IF disables these interrupts.

Architecture

8086 registers categorized into 4 groups





Sl.No.	Type	Register width	Name of register
1	General purpose register	16 bit	AX, BX, CX, DX
		8 bit	AL, AH, BL, BH, CL, CH, DL, DH
2	Pointer register	16 bit	SP, BP
3	Index register	16 bit	SI, DI
4	Instruction Pointer	16 bit	IP
5	Segment register	16 bit	CS, DS, SS, ES
6	Flag (PSW)	16 bit	Flag register

Architecture

Registers and Special Functions

Register	Name of the Register	Special Function
AX	16-bit Accumulator	Stores the 16-bit results of arithmetic and logic operations
AL	8-bit Accumulator	Stores the 8-bit results of arithmetic and logic operations
BX	Base register	Used to hold base value in base addressing mode to access memory data
CX	Count Register	Used to hold the count value in SHIFT, ROTATE and LOOP instructions
DX	Data Register	Used to hold data for multiplication and division operations
SP	Stack Pointer	Used to hold the offset address of top stack memory
BP	Base Pointer	Used to hold the base value in base addressing using SS register to access data from stack memory
SI	Source Index	Used to hold index value of source operand (data) for string instructions
DI	Data Index	Used to hold the index value of destination operand (data) for string operations